

# **Trust Model**

## **Supplements and Taxonomy**

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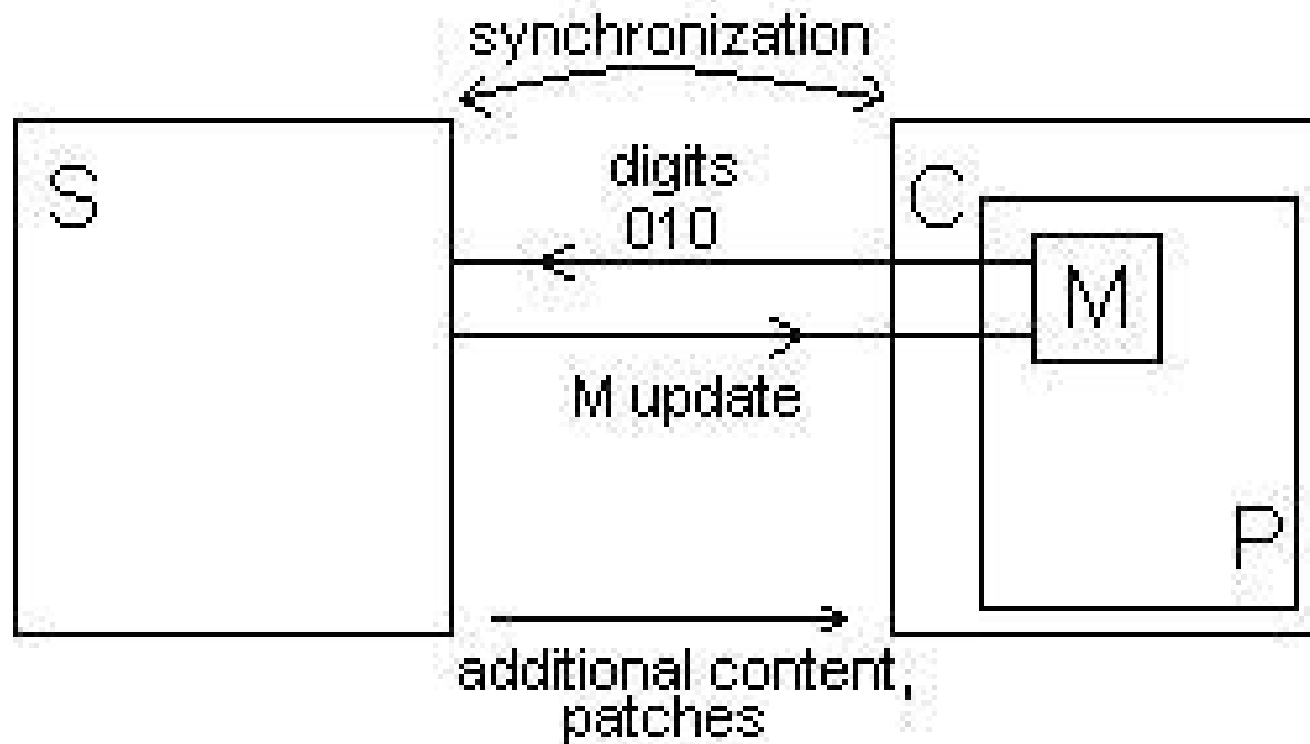
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# Introduction

- Trust Model





## Distribution of verification functions between client program and trusted server (1/2)

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- It is possible to execute some verification functions within the trusted server
- Client monitor has two responsibilities:
  - To **gather** needed data about a program and its state
  - To **execute** a part of verification functions
- Trusted server has two responsibilities:
  - To **interpret** all received signatures to authenticate the client program
  - To get and **extract** gathered data from client to execute some verification functions itself



## Distribution of verification functions between client program and trusted server (2/2)

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- Problem: what verifications should be executed on
  - trusted server?
  - client monitor?
  - both?
- To approach the problem solution, we may take into account the following verification properties:
  - computational complexity
  - relative importance
  - size of input data and their structure



## Monitor update (1/2)

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- We need to construct an infinite sequence of monitors  $M_0, M_1, M_2, \dots$  so that
  - $\{\exists \text{const1} > 0: \forall k \geq 0 \text{ time\_of\_breaking}(M_k) > \text{const1}\}$   
– monitor level
  - Or
  - $\forall i \geq 0$   
 $\{\exists \text{const2} > 0: \forall k \geq 0 \text{ time\_of\_breaking}(VFi(M_k)) > \text{const2}\}$   
– verification function level; where  $VFi(M_k)$  is some verification number  $i$  of  $M_k$  monitor
- So our aim is to design appropriate sequences of verification functions and also data gather functions



## Monitor update (2/2)

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- Each verification function may defer from its previous versions by some
  - **regular modifications**, e.g.
    - modification of secret key
    - hash function modification
    - sequence and amount of gathered data
    - sequence of passing Control Flow Graph
  - and **qualitative modifications**, such as
    - adding a new type of verifications and gathering functions
    - deleting some existing out-of-date verifications and gathering functions



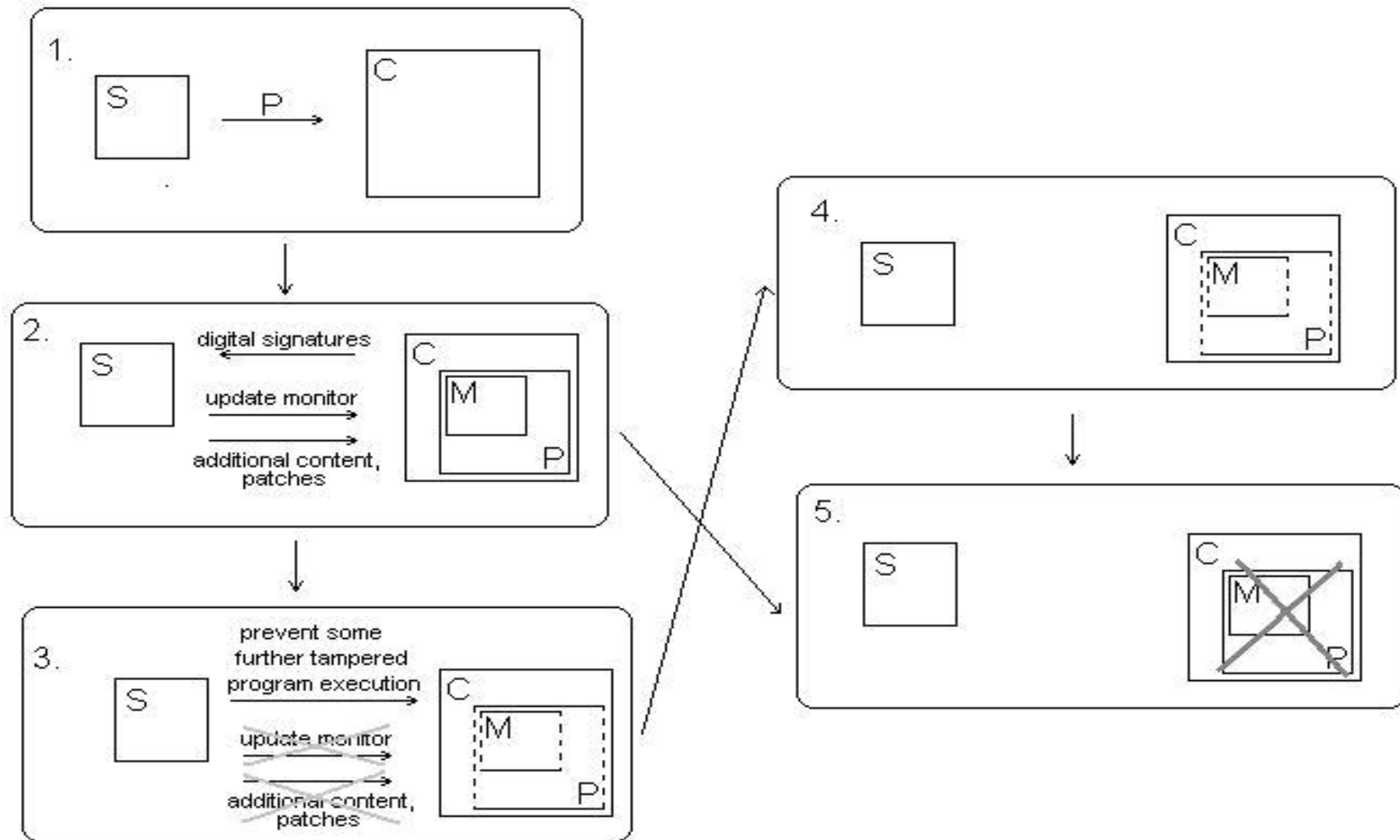
## Actions executed by trusted server and/or by client when detecting a client program tampering

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- Suspend to send additional content, patches and system data
- Suspend the monitor update
- Send to the client certain commands that execute some irreversible program modifications preventing the program work:
  - destruction of the program parts which are difficult to restore;
  - definite or arbitrary modification of program;
  - definite or arbitrary modification or removal of data and files used by program.



# Client program life cycle





# Trust model taxonomy (1/4)

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- In the current version of **taxonomy** there are three high-level notions:
  - Client
  - Trusted server
  - Data streams between them



# Trust model taxonomy (2/4)

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- Client
  - Target program
    - Monitor
      - Verification functions
        - Static verifications
        - Integrity verification of program code; implementation: check-sums and hash-functions
        - Dynamic verification
        - Calculation results correctness
        - Assert-instructions
        - Environment verification: HW/OS
        - CFG correctness verification. Notions: basic blocks, marks on them, CFG paths
      - Functions of **data gathering** for verification on server
      - **Reaction** on illegal modifications
        - Destruction of difficult restoring parts of program
        - Definite or arbitrary program modification
        - Definite or arbitrary program modification or removal of data and files used by program
    - Proper client program



# Trust model taxonomy (3/4)

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- Server
  - Signature interpreting component
  - Monitor factory: mechanism, routinely generating a new monitor version
    - Regular modification
      - Secret key modification
      - Hash function modification
      - Modification of sequence and amount of gathered data
      - Modification of sequence of passing vertex of Control Flow Graph
    - Qualitative modification
      - Addition of new verifications type
      - Addition of new gathering functions type
      - Removal out-of-date verifications and gathering functions (for time optimization)
  - Manager, determining “intension” of entire protection functioning:
    - monitor update periodicity
    - periodicity of receiving required signatures
    - working life of signatures and others
    - current state of client (tampered/not tampered)
  - Verification functions



# Trust model taxonomy (4/4)

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- Data streams
  - Streams from client to server
    - Authentication signatures
    - Data for verification on server
  - Streams from server to client
    - Monitor updates
    - Additional content and patches
  - System data, organizing a single physical data stream between client and server and synchronizing client with server



# Conclusion

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- So, we have touched several **unsolved problems**:
  - **How to distribute** verification functions between client program and trusted server?
  - What methods may be used to **construct verification sequences** for monitor update?
  - The problem on **acceptability** of additional preventing actions when detecting a client program tampering.