

# *Hypothetical* **Trust and Attack Models**

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# Outline

- Architecture of the software-only solution.
- Sources of trust.
- Possible attacks.
- Open issues.

# The remote entrusting problem

*Remote software authentication*: ensuring a trusted machine (server) that an untrusted host (client) is running a “healthy” version of a program  $P$ :

- The program is unadulterated.
- It is executed on top of unadulterated HW/SW.
- The execution process is not manipulated externally.

The distinctive feature of *remote* entrusting is that the authenticated software needs to communicate over the network with the trusted machine to work properly.

# Entrusting remote applications

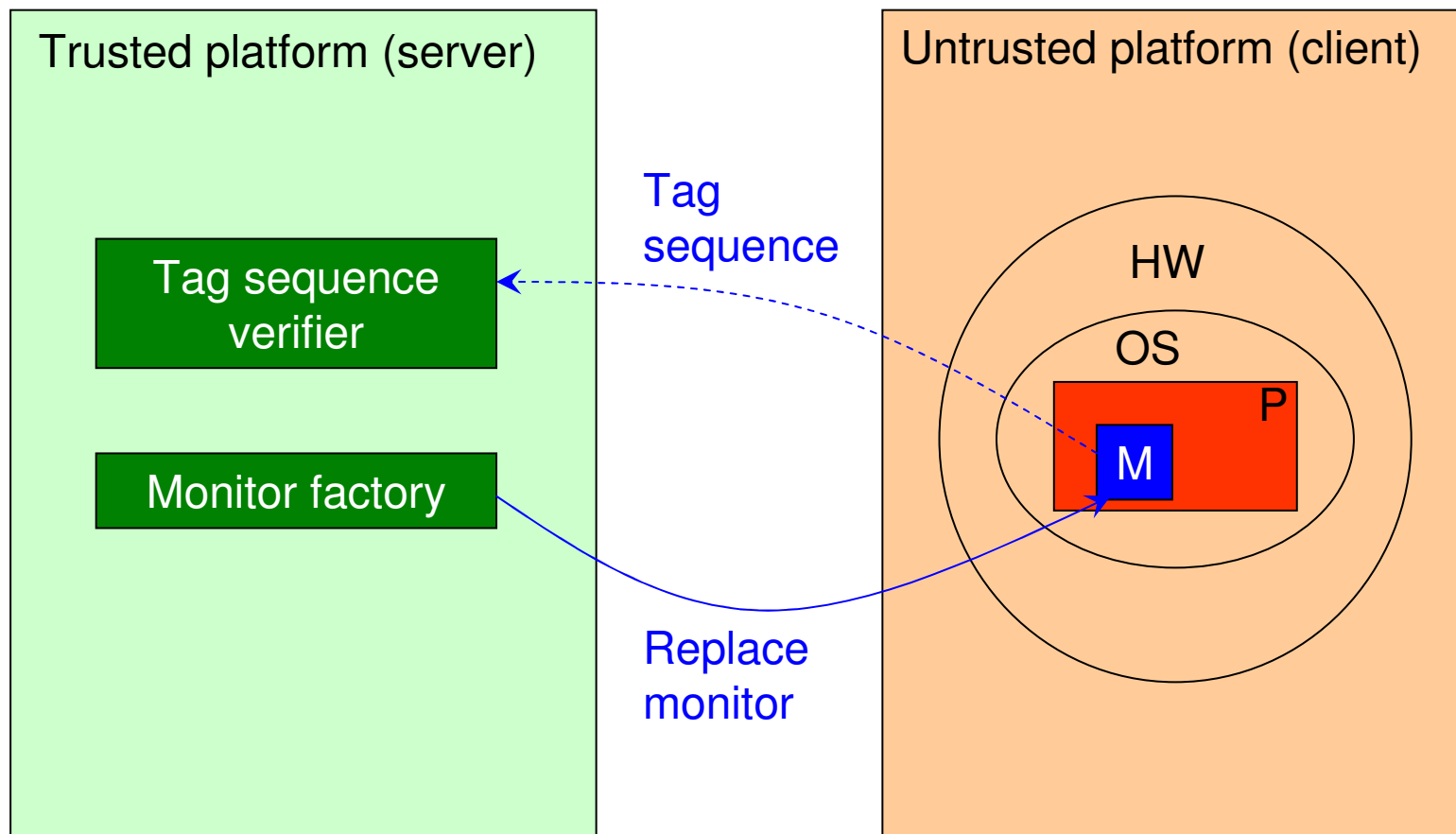
- An increasing number of applications depend on services provided over the network.
- Often service providers need to “entrust” their clients and assume they do not act maliciously.
- This is critical for applications involving:
  - Trading (eCommerce), bidding, gaming.
  - eGovernment, eVoting.
  - Distributed (Grid) computing.
  - Protocol implementations.
  - Content and data protection.



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# Reference architecture



# Sources of trust

# Authenticity verification

The monitor  $M$  should verify:

- Text and data segments of  $P$  as loaded in memory.
- Libraries used by  $P$ .
- The execution environment (HW, OS, execution process, etc.).
- Results of specific computations or assertions.

# Tag sequence generation

The monitor  $M$  sends the server an authenticity tag sequence as evidence of healthy execution:

- Tags have limited time validity.
- A secret key, hidden into  $M$  itself, is used to generate them.
- If no tag or an incorrect tag is received by the server, the client is considered untrusted and no service is delivered to it.



# Replacement

To give attackers a limited time to succeed, the monitor  $M$  is periodically replaced:

- The duration depends on the estimated reverse engineering complexity, assuming humans are necessarily involved in the process.
- The monitor factory should generate highly independent monitors.

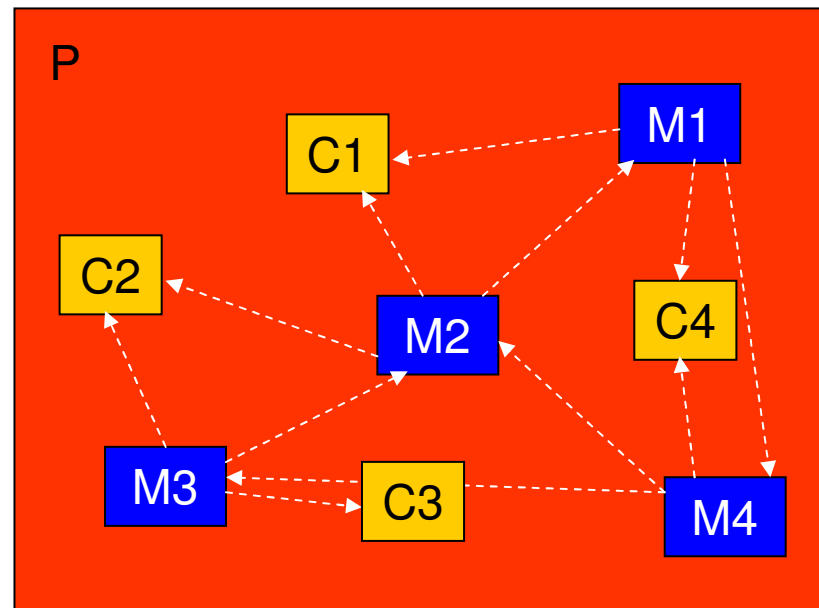
# Code obfuscation

To increase the resistance to reverse engineering, the code is obfuscated:

- Opaque predicates based on conditions that are hard to analyze statically (e.g., involving pointer structures) could be used.

# Additional sources of trust

- **Self checking monitor:**  $M$  checks itself before checking  $P$ .
- **Network of trust:**
- **Tags include data verified by server:**  
authenticity verification is no longer local to  $M$
- **Server sends challenge  $C$  to client:** tag generation and authenticity verification depend<sub>11</sub> on  $C$ .



# Attacks

# Assumptions on attackers

A malicious user can:

- Install any software on the client.
- Read and write memory locations, processor registers and files.
- Observe and modify the network traffic.
- Modify  $P$  and  $M$ , both on disk and in memory.
- Use any available code analysis tool.
- Take advantage of tracers, emulators and debuggers.
- Tamper with libraries, operating system and hardware.

A malicious user cannot:

- Access and tamper with the trusted server.
- Know the software/hardware configuration of the server.



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# Classes of attacks

1. Reverse engineering attack.
2. Execution environment attack.
3. Cloning attack.
4. Differential analysis attack

# Reverse engineering attacks

Important functionalities and data structures are located and altered maliciously in  $P$  and  $M$ :

- Tag sequence generator.
- Authenticity checking functions.
- Secret keys.
- Input data (e.g., passed to checking functions).
- Output data (e.g., returned by checking functions).

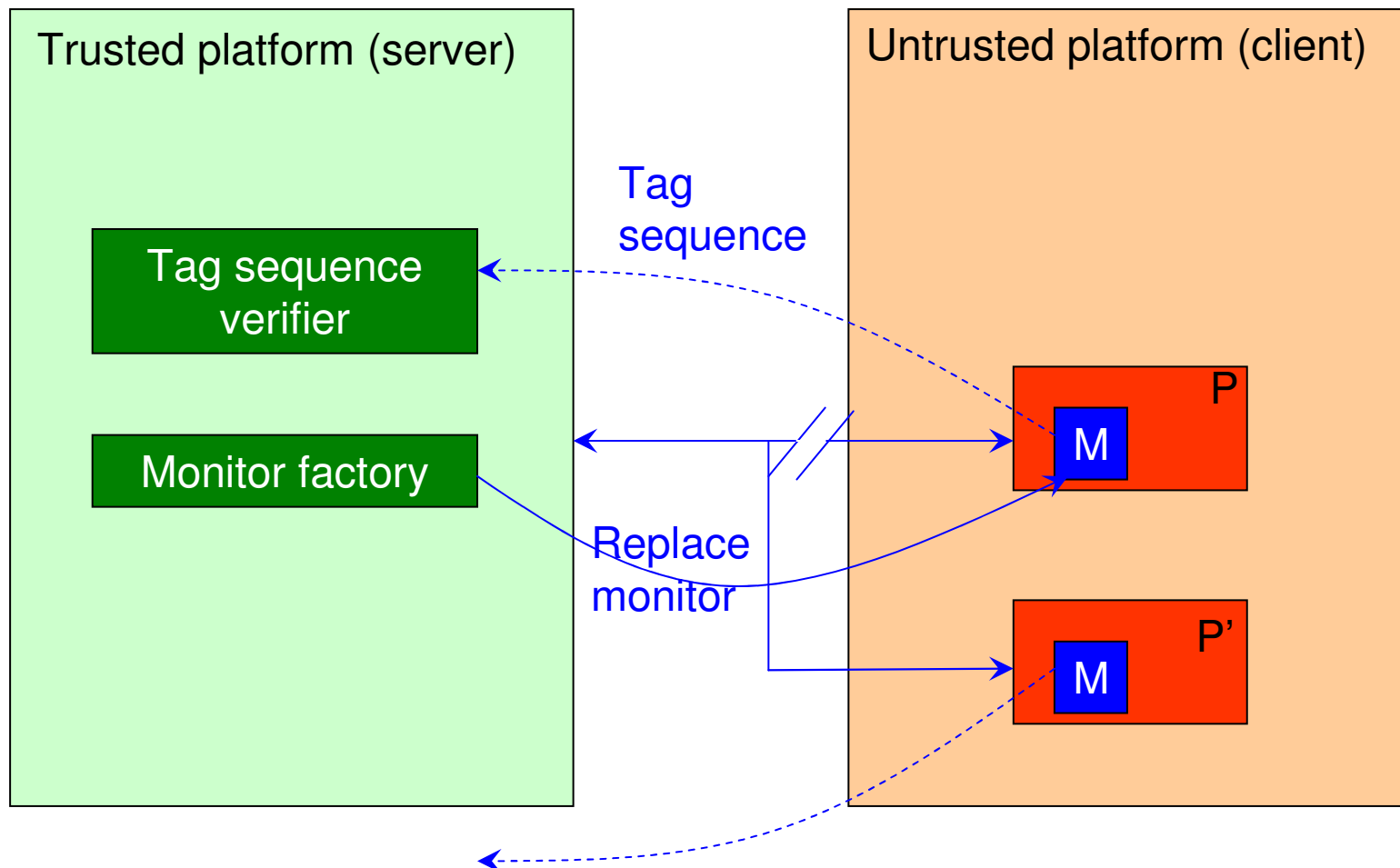
# Execution environment attacks

$P$  is run on an emulator, in debug mode or is interpreted by an adulterated virtual machine:

- Memory locations, call stack, program counter and parameters can be altered dynamically.
- Dynamic libraries can be altered maliciously.
- Input and output values can be replaced on-the-fly.



# Cloning attack

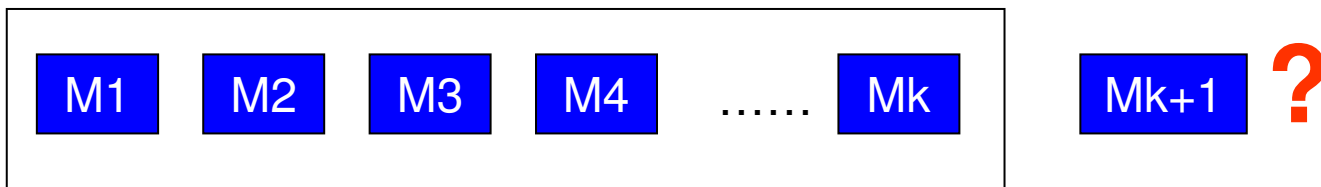


This attack is ineffective if tag sequence includes computation data.

# Differential analysis attack

The attacker gathers information about  $M$  by comparing the sequence of monitors delivered by the monitor factory in the past:

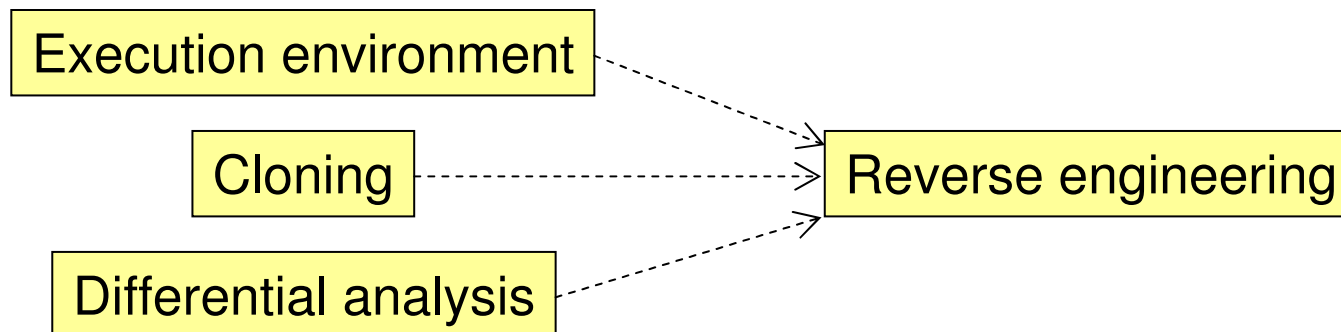
- If the strategy used by the monitor factory is even only partially understood, the time necessary to break new monitors might be reduced, eventually allowing the attacker to break a yet valid monitor.



# Dependencies among attacks

Execution environment, cloning and differential analysis attacks are effective only if combined with reverse engineering, which provides information on:

- Where functionalities and data structures are implemented.
- When to intercept the execution.
- How to alter it.





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# Analysis of attack resistance

|     | A1 | A2                                               | A3 | A4 | A5 | A6 | A7 | A8 | A9              | A10 | A11 | A12 |    |  |
|-----|----|--------------------------------------------------|----|----|----|----|----|----|-----------------|-----|-----|-----|----|--|
| T1  | x  | T8: Monitor replacement                          |    |    |    |    |    |    |                 |     |     |     |    |  |
| T2  |    | T1: M checks P text and data segment             |    |    |    |    |    |    |                 |     |     |     |    |  |
| T3  |    | T2: M self checks itself before checking P       |    |    |    |    |    |    | x               |     |     |     |    |  |
| T4  |    | T3: M checks libraries used by P                 |    |    |    |    |    |    |                 | x   |     |     |    |  |
| T5  |    | T4: M checks execution environment               |    |    |    |    |    |    | ge from server) |     |     |     |    |  |
| T6  | x  | T5: M checks OS and HW on env.                   |    |    |    |    |    |    |                 |     |     |     |    |  |
| T7  |    | T6: M checks results of computation              |    |    |    |    |    |    |                 |     |     |     |    |  |
| T8  |    | T7: Secret key used to generate tag sequence     |    |    |    |    |    |    |                 |     |     |     | ks |  |
| T9  | x  | A3: Replace tag sequence generator               |    |    |    |    |    |    |                 |     |     |     | x  |  |
| T10 |    | A5: Modify input before call on env.             |    |    |    |    |    |    |                 |     |     |     |    |  |
| T11 |    | A6: Modify output before return M/P              |    |    |    |    |    |    |                 |     |     |     |    |  |
| T12 |    | A7: Modify output before return on env.          |    |    |    |    |    |    |                 |     |     |     |    |  |
|     |    | A8, A9: Replace HW/OS; replace dynamic libraries |    |    |    |    |    |    |                 |     |     |     |    |  |
|     |    | A10: Tampered execution (debug mode)             |    |    |    |    |    |    |                 |     |     |     |    |  |

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# Open issues (1)

- Protection against reverse engineering:
  - Code obfuscation metrics
  - Beyond code obfuscation?
- Monitor factory:
  - How to construct independent monitors automatically?
  - What is the minimum set of functionalities of a monitor?
- Network of trust:
  - What replacement frequency can we achieve with it?
- Tags include output data:
  - How to avoid/minimize replication of the computation on the server?

## Open issues (2)

- Server sends challenge to client:
  - How does it complement/substitute monitor replacement?
- Tag sequence generator:
  - How to hide secret key into the monitor's code?
- Execution environment:
  - How to check sanity of HW and execution mode?
  - How to recognize parallel execution of clone process?

# Conclusions

- The reference architecture alone is very fragile and should be augmented with:
  - Network of trust
  - Challenge from server
  - Data embedded into tag sequence
- Replacement and monitor factory are at the core of remote entrusting
  - Independent monitor generation is crucial



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# Proposal for a new reference architecture

