

SPIIRAS Team in RE-TRUST: team background and preliminary analysis of tasks to be solved

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 - Modeling and simulation of computer attacks
 - Security analysis of computer networks
 - Intrusion detection
 - Deception systems, honeynets
 - Modeling and simulation of cyberwarfare
 - Security policy specification and checking
 - Security protocols analysis



1. St. Petersburg Institute for Informatics and Automation of the Russian Academy of Sciences (SPIIRAS)



Short Profile of St. Petersburg Institute for Informatics and Automation (SPIIRAS)

- Founded in 1978
- The Russian Academy of Sciences Institute operating in the North-West of Russia in Information Technologies
- Personnel 203
- SPIIRAS is a competence center in the area of advanced information technologies
- Experienced in collaboration with EC Countries



SPIIRAS Structure



SPIIRAS General Research Directions

- Development of Information and Control Systems for Real Time Signal Processing
- Fundamentals of Information Processes in Complex (Socio-, Eco-, Bio-,Geo-, etc) Systems
- Theoretic Basics in Developing Information Technologies for Research Automation, Control, Manufacturing, and Intelligent Systems

Research Directions of Computer Security Research Group

- Computer security, including security policy management, access control, authentication, network security analysis, intrusion detection, firewalls, deception systems, virus protection, verification of security systems, modeling, simulation and visualization technologies for counteraction to cyber terrorism;
- Artificial intelligence, including multi-agent frameworks and systems, soft and evolutionary computing, machine learning, data mining, data and information fusion;
- **Telecommunication**, including decision making and planning for telecommunication systems.



2. International collaboration, projects and accomplishments of Computer Security Research Group

International Projects

- Air Force Research Laboratory/ Information Directorate (European Office of Aerospace Research and Development) (1999-2003 - 3 computer security projects)
 - "Agent-Based Model of Information Security System: Architecture and Formal Framework for Coordinated Intelligent Agents Behavior Specification" (1999-2001)
 - "Formal Grammar-Based Approach and Tool for Simulation Attacks against Computer Network" (2001-2003)
 - "Mathematical Foundations, Architecture and Principles of Implementation of Multi-Agent Learning Components for Attack Detection Computer Networks" (2001-2003)
- INTEL "Network traffic preprocessing algorithms" (2004-2005)
- Fraunhofer First (Germany) "Intrusion detection learning systems" (MIND) (2004-2006)
- FP6 (EU Project) "Security policy specification, checking and deployment" (POSITIF) (2004-2007)
- ..
- FP6 (EU Project) "Remote EnTrusting by RUn-time Software auThentication" (RE-TRUST) (2006-2009)

Russian Grants and Projects

- Government Budget Projects:
 - Models and methods of developing secure computer systems (2006-2008)
- Grants of Russian Foundation for Basic Research:
 - Mathematical models of information security assurance in computer networks based on MAS technology and its experimental evaluation (2001-2003)
 - Modeling and simulation of cyber warfare (2004-2006)
- Projects from Department of Information Technology and Computer Systems of the Russian Academy of Sciences:
 - Agent-based stochastic modeling and simulation of adversarial competition of teams in the Internet environment (2003-2004)
 - Mathematical models of active audit of computer network vulnerabilities, intrusion detection and response: Multi-agent approach (2003-2006)
- Projects from Government Institutions:
 - Models and prototypes of honeynets (2003-2004)
 - Monitoring of user activity in computer networks (2004)
 - **RE-TRUST Kick-off Workshop, September 18-19, 2006**

International Computer Security Conferences Organized

 Mathematical methods, models and architectures for computer network security (MMM-ACNS): 2001, 2003, 2005, 2007



NATSON





The objectives of MMM-ACNS Workshops are to bring together leading researchers from academia and governmental organizations as well as practitioners in the area of computer networks and information security, facilitating personal interactions and discussions on various aspects of information technologies in conjunction with computer network and information security problems arising in largescale computer networks engaged in information storing, transmitting, and processing.

Recent accomplishments (results) in 2005-2006 (1)

- The theoretical basis, algorithms and software implementation of agentoriented modeling and simulation of antagonistic counteraction of malefactors and computer network security components.
 - The principles of construction, the structure and a fragment of distributed ontology-based knowledge base for modeling and simulation of protection mechanisms in an antagonistic environment.
 - The formal models of agents-malefactors, security agents and computer network under defense.
 - The software environment based on imitation of computer attacks and protection mechanisms at a network packets level using OMNeT ++ INET Framework.
 - The attack and security agents as compound modules containing simple modules, responsible for functioning of various network protocols, and agent kernel that controls these modules.
 - Different experiments with this environment were fulfilled on an example of simulation of DDoS attacks and particular defense mechanisms. These experiments included the investigation of attack scenarios and protection mechanisms for the networks with different structures and security policies.
 - The received results are directed on investigation of various aspects of antagonistic interactions of agent teams in the Internet and development of recommendations on design and implementation of advanced security systems.

Recent accomplishments (results) in 2005-2006 (2)

- The theoretical basis and operation algorithms of deception systems (DS).
 - These systems represent hardware-software tools for information protection that are based on the technology of "traps" and false targets. In particular, we developed the requirements to DS, the generalized architecture of multi-agent DS, the generalized models and algorithms of disguised counteraction to remote non-authorized access to information resources, including the models of malefactor detection and readdressing of non-authorized request to false components, determining the malefactor plan (strategy), generating a plan of false components operation, etc.
 - The offered approach is based on simulation of information systems components and on using three levels of malefactor deception: (1) a network segment level – the whole network segment is emulated; (2) a host level – among working servers the bait-host is used; (3) a services and applications level –the programs emulating services and applications are applied on servers.
 - The deception software system was implemented.
 - We fulfilled a set of experiments on investigating basic deception functions at realization of different attacks. These experiments are executed on several different scenarios determined according to various attack types.

Recent accomplishments (results) in 2005-2006 (3)

- Models, techniques and prototypes for active security analysis of computer networks.
 - The approach is based on automatic generation and fulfillment of distributed attack scripts taking into account a variety of goals and knowledge levels of malefactors, and intended for implementation at various stages of computer network life cycle including design and exploitation stages.
 - The offered approach is based on application of a set of models (including models of malefactor, attack scripts generation, security level evaluation, computer network, etc.) using expert knowledge.
 - Functioning of security analysis system based on the approach suggested is resulted in determined vulnerabilities, the traces (graphs) of possible multistage attacks, "bottlenecks" (main "holes") in a computer network on which these attacks are based, and also various security metrics which can be used for evaluating a security level of a computer network and its components, and also for comparisons of various network configurations and security policies.
 - These results provide the development of justified recommendations on elimination of revealed "bottlenecks" and on amplification of system security.

Recent accomplishments (results) in 2005-2006 (4)

- The generalized architecture, particular models and prototypes of components for verifying security policies of computer networks were analyzed and developed.
 - The mechanisms for operation with policies of three levels were offered:
 - (1) the top level, approximated to the user requirement language,
 - (2) the intermediate level, classifying rules according to several policy categories, and
 - (3) the bottom level, describing policies in the format of Common Information Model (CIM).
 - We developed and implemented the research prototypes of the verification manager that handles the process of verification, and different verification modules:
 - (1) based on Event Calculus,
 - (2) based on Model Checking,
 - (3) Specialized modules.

Some Recent research papers published or accepted for publishing in 2006 (1)

- Kotenko I., Stepashkin M., Ulanov A. Agent-based modeling and simulation of malefactors' attacks against computer networks. Security and Embedded Systems. D.N.Serpanos, R.Giladi (Eds.). IOS Press. 2006.
- Kotenko I.V., Tishkov A.V., Chervatuk O.V. Architecture and Models for Security Policy Verification. Mathematics and Security of Information Technologies. Amsterdam, 2006.
- Kotenko I., Ulanov A. Agent-based Simulation of Distributed Defense against Computer Network Attacks. Proceedings of 20th European Conference on Modelling and Simulation (ECMS 2006). Bonn. Germany. May 28th - 31st, 2006. P.560-565.
- Kotenko I., Ulanov A. Antagonistic Agents in the Internet: Computer Network Warfare Simulation. The 9th International Conference on Information Fusion. Florence (Italy), 10-13 July, 2006.
- Kotenko I., Stepashkin M. Network Security Evaluation based on Simulation of Malefactor's Behavior. SECRYPT International Conference on Security and Cryptography. International Joint Conference on e-Business and Telecommunications. ICETE 2006. Setubal, Portugal. 7-10 August 2006.

Some Recent research papers published or accepted for publishing in 2006 (2)

- Kotenko I., Ulanov A. Agent-based modeling and simulation of network softbots' competition. The Joint Conference on Knowledge-Based Software Engineering (JCKBSE'06). Tallinn, Estonia. August 28-31. 2006.
- Kotenko I., Ulanov A. Simulation of Internet DDoS Attacks and Defense. 9th Information Security Conference. ISC 2006. Samos, Greece. August 30 -September 2, 2006. Lecture Notes in Computer Science, Vol. 4176, 2006.
- Kotenko I., Stepashkin M. Analyzing network security using malefactor action graphs. International Journal of Computer Science & Network Security, 2006.
- Kotenko I., Stepashkin M. Attack Graph based Evaluation of Network Security. The 10th IFIP Conference on Communications and Multimedia Security. CMS'2006. Heraklion, Greece. 19 - 21 October 2006. Lecture Notes in Computer Science, Vol. 4237, 2006.
- Kotenko I., Ulanov A. Agent Teams in Cyberspace: Security Guards in the Global Internet // International Conference on CYBERWORLDS. CW2006. Lausanne, Switzerland, November 28-30, 2006. Proceedings. IEEE Computer Society, 2006.



3. Tasks in RE-TRUST

Main Tasks in RE-TRUST

- Task T4.1 Trust analysis of SW-based method
- Task T4.5 Remote entrusting and Internet secure protocols
- T4.5.1 Analysis of integration of remote entrusting with existing Internet security protocols
- T4.5.2 Integration and analysis of secure protocols to support remote entrusting methods.



4. Computer Security Research Group Research: General view

Examples of related developed solutions (mainly in network security analysis):

- Modeling and simulation of computer attacks
- Modeling and simulation of cyberwarfare
- Security analysis of computer networks
- Intrusion detection
- Deception systems, honeynets
- Security policy specification and checking
- Security protocols analysis

Security Evaluation Areas

- Impact assessment for determining how security measures affect system and application properties (performance, reliability, etc.) [D.Nicol, S.Smith, M.Zhao-04; S.Kent, C.Lynn, K.Seo-00 (Secure BGP); M.Zhao, S.Smith, D.Nicol-05; etc.]
- Emulation, in which real and virtual worlds are combined to study the interaction between malware and systems, and probe for new system weaknesses [G.Bakos, V.Berk-02 (Worm activity by metering ICMP); M. Liljenstam et al-03 (Simulating worm traffic); etc.]
- Cyber attack exercises and training scenarios [M. Liljenstam et al-05 (RINSE); B. Brown et al-03; etc.]

. . .

 Risk assessment based on known vulnerabilities, exploits, attack capabilities, and system configuration [R. Ortalo, Y.Deswarte, M.Kaaniche-99; Sheyner et al-02; V.Gorodetski, I.Kotenko-02 (Attack Simulator); B.Madam, K.Goseva-Popstojanova-02; etc.]

Works on Security Evaluation

- Methodology and software tools for testing IDSs ([Puketza *et al*-96], [Puketza *et al*-97], [Debar *et al*-98], [Alessandri *et al*-01], [McHugh-00]);
- Evaluations of IDSs of MIT ([Lippmann et al-98, 00, 02]);
- Real-time test bed of AFRL [Durst et al-00];
- Dependability models for evaluation security [Nicol et al-04];
- Penetration testing of formal models of networks for estimating security metrics [Sheyner *et al*-02];
- Model checking for analysis of network vulnerabilities [Ritchey, Ammann-00];
- Global metrics for analyzing the effects of complex network faults and attacks [Hariri *et al*-03];
- Natural-deduction for automatic generation and analysis of attacks against IDS [Rubin *et al*-04];
- Knowledge-based network risk assessment [Shepard et al-05], etc.

Works directly coupled with Attack Modeling and Simulation

- Using Colored Petri Nets [Kumar et al-94];
- State transition analysis technique [Iglun *et al*-95], [Kemmerer *et al*-98];
- Conceptual models of computer penetration ([Cohen-99],[Stewart-99]);
- Descriptive models of attackers [Yuill et al-00];
- "Tree"-based models of attacks ([Moore et al-01], [Dawkins et al-02]);
- Modeling survivability of networked systems [Moitra et al-01];
- Object-oriented Discrete Event Simulation [Chi et al-01];
- Situation calculus and goal-directed procedure invocation [Goldman-02];
- Using and building attack graphs for vulnerability analysis ([Swiler *et al*-01], [Ortalo *et al*-01], [Sheyner *et al*-02], [Jha *et al*-02]);
- Game-theoretic models [Lye and Wing-03];
- Multi-stage attack analysis [Dawkins, Hale-04];
- Modeling and inference of attacker intent, objectives, and strategies [Liu, Zang-05]; etc.

Security Analysis

- Model system
- 2 Model adversary
- 3 Identify security properties
- 4 See if properties preserved under attack
- Result
 - Under given assumptions about system, no attack of a certain form will destroy specified properties
 - There is no "absolute" security

/Vitaly Shmatikov/

Fundamental Tradeoff

- Formal models are abstract and greatly simplified
 - Components modeled as finite-state machines
 - Security functions modeled as abstract data types
 - Security property stated as unreachability of "bad" state
- Formal models are tractable...
 - Lots of verification methods, many automated
- ...but not necessarily sound
 - Proofs in the abstract model are subject to simplifying assumptions which ignore some of attacker's capabilities
- Attack in the formal model implies actual attack

/Vitaly Shmatikov/





• Modeling and simulation of computer attacks



Basic Components of Attack Model

- 1. **Ontology of the Problem** "*Attacks against Computer Network*" : structure of the basic malefactors' intentions and actions.
- 2. Basic malefactors' intentions and attack task specification.
- 3. Formal grammar-based framework for specification of attack development.
- 4. Formal scenarios of a **representative multitude of attacks** and their development in time.
- 5. Formal model of the attacked computer network.
- 6. Model of **interaction** of malefactor's activity and victim computer network.



Basic malefactors' intentions

Intention-centric approach to the specification of malefactor's activity: basic notions of the domain correspond to the malefactor intentions and all other notions are structured according to the structure of intentions.

List of Basic Classes of High-level Malefactor's Intentions

- R Reconnaissance:
 - *IH Identification* of the running *Hosts*
 - *IS Identification* of the host *Services*
 - *IO Identification* of the host *Operating* system
 - *CI Collection* of additional *Information*
 - *RE* shared *Resource Enumeration*
 - *UE Users* and groups *Enumeration*
 - ABE Applications and Banners Enumeration

- *I Implantation* and threat realization:
 - GAR Getting Access to Resources of the host
 - *EP Escalating Privilege* with regard to the host resources
 - GAD Gaining Additional Data needed for further threat realization
 - TR Threat Realization
 - *CD Confidentiality Destruction*
 - *ID Integrity Destruction*
 - DOS Violation of resource availability (Denial of Service)
 - CT Covering Tracks CBD – Creating Back Doors

Attack task specification

ecify the At	Lack		
Intention I IH IIH IIH IIH III IIS III IIH III III I	Description Identification of Hosts Identification of Services Identification of Operating system Shared Resource Enumeration Applications and Banners Enumeration Getting Access to Resources of the host Escalating Privilege with regard to the ho Confidentiality Violation Realization Integrity Violation Realization Creating Back Doors guration ass 161 43 201 148 C D:\HACKER\files\passwd.txt Net IP I 32.168.130.0 210.122.25.0 E Known Information	Save preceding attack realization Generate attacks on net protocol level s st Name Host IP 210.122.25.12 210.122.25.12 210.122.25.12 210.122.25.22 E 192.168.130.137 Show All Hosts Advanced	 Main elements of attack specification: 1) Malefactor's intention (1-12); 2) Address of the attacked host or network; 3) Available information about attacked host; 4) Attack object (file name, user account, resource, etc.);

Formal framework for specification of attacks

Formal grammar: $G_i = \langle V_N, V_T, S, P, A \rangle$,

where G_i – formal grammar name (it coincides with the name of attack and the name of its axiom);

- V_N the set of non-terminal symbols; V_T the set of terminal symbols; $S \in V_N$ formal grammar axiom;
- **P** the set of productions which look like follows:

(U) $X \rightarrow \alpha$ (Prob),

- where $X \in V_N$, $\alpha \in (V_T \cup V_N)^*$, U precondition of the production application; **Prob** probability of the production application;
- A the set of attributes and their dependencies (functions having attributes as variables).

Implementation Issue: State Machine-based Representation of Attack Generation



Implementation Issue: State Machine-based Representation of Attack Generation



"Reconnaissance" Attack Generation *IR1* – Intermediate state

- IH Identification of running Hosts
- **IS** Identification of Services
- **IO** Identification of OS
- CI Collection of Information
- **RE** Resource Enumeration
- UE Users and Groups Enumeration
- **ABE** Applications and Banners Enumeration
User Interface with Network Model





On-line Visualization of an Attack Development on Micro-Level

Shortcut to PORTAL.BAT	- 🗆 🗙 Shortcut to PORTAL.BAT
tarting scanports v.1.0. TCP scanning by using SYN messages. ttackID: SS	3. 192.168.130.136.1050->192.168.130.135.81 TCP RST ACK (seq: 12f799 ack: 1)
elected device: Realtek 8139-series PCI NIC	SYN flooding v.1.0 Starting
. 192.168.130.136.1050->192.168.130.135.21 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.21->192.168.130.136.1050 TCP SYN ACK (seq: 8b6feee8 ack: 12f799) ort 21 is seems to be OPEN. . 192.168.130.136.1050->192.168.130.135.21 TCP RST ACK (seq: 12f799 ack: 8b6feee9)	102 160 120 1E 1020 \102 160 120 12E 21 TCP CUN (opg: 24270 pak; 0)
. 192.168.130.136.1050->192.168.130.135.79 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.79->192.168.130.136.1050 TCP RST ACK (seq: 0 ack: 12f799) ort 79 is seems to be CLOSED. . 192.168.130.136.1050->192.168.130.135.79 TCP RST ACK (seq: 12f799 ack: 1)	192.168.128.15.1030->192.168.130.135.21 TCP SYN (seq: 15e0a ack: 0) 192.168.128.15.1031->192.168.130.135.21 TCP SYN (seq: 15F90 ack: 0) 192.168.128.15.1032->192.168.130.135.21 TCP SYN (seq: 214b2 ack: 0) 192.168.128.15.1033->192.168.130.135.21 TCP SYN (seq: 23451 ack: 0) 192.168.128.15.1033->192.168.130.135.21 TCP SYN (seq: 23451 ack: 0) 192.168.128.15.1033->192.168.130.135.21 TCP SYN (seq: 23451 ack: 0)
. 192.168.130.136.1050->192.168.130.135.80 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.80->192.168.130.136.1050 TCP SYN ACK (seq: 8b788c3f ack: 12f799) ort 80 is seems to be OPEN. . 192.168.130.136.1050->192.168.130.135.80 TCP RST ACK (seq: 12f799 ack: 8b788c40)	192.168.128.15.1035->192.168.130.135.21 TCP SYN (seq: 25a62 ack: 0) 192.168.128.15.1036->192.168.130.135.21 TCP SYN (seq: b8ab ack: 0) 192.168.128.15.1037->192.168.130.135.21 TCP SYN (seq: 2436 ack: 0) 192.168.128.15.1038->192.168.130.135.21 TCP SYN (seq: 36f1 ack: 0) 192.168.128.15.1039->192.168.130.135.21 TCP SYN (seq: 36f1 ack: 0)
. 192.168.130.136.1050->192.168.130.135.81 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.81->192.168.130.136.1050 TCP RST ACK (seq: 0 ack: 12f799) ort 81 is seems to be CLOSED. . 192.168.130.136.1050->192.168.130.135.81 TCP RST ACK (seq: 12f799 ack: 1)	192.168.128.15.1040-)192.168.130.135.21 TCP SVN (seq: 31a1 ack: 0) 192.168.128.15.1041-)192.168.130.135.21 TCP SVN (seq: 31a1 ack: 0) 192.168.128.15.1041-)192.168.130.135.21 TCP SVN (seq: 1a20c ack: 0) 192.168.128.15.1043-)192.168.130.135.21 TCP SVN (seq: 7a73 ack: 0) 192.168.128.15.1043-)192.168.130.135.21 TCP SVN (seq: 202cc ack: 0) 192.168.128.15.1043-)192.168.130.135.21 TCP SVN (seq: 202cc ack: 0)
tarting scanports v.1.0. TCP scanning by using SYN messages. ttackID: HS	192.168.128.15.1045->192.168.130.135.21 TCP SYN (seq: 18f51 ack: 0) 192.168.128.15.1046->192.168.130.135.21 TCP SYN (seq: 13f51 ack: 0) 192.168.128.15.1046->192.168.130.135.21 TCP SYN (seq: 13fc5 ack: 0) 192.168.128.15.1046->192.168.130.135.21 TCP SYN (seq: 2472 ack: 0)
elected device: Realtek 8139-series PCI NIC	192.168.128.15.1048-)192.168.130.135.21 TCP SYN (seq: 14501 ack: 0) 192.168.128.15.1049->192.168.130.135.21 TCP SYN (seq: 3d63 ack: 0)
. 192.168.130.136.1050->192.168.130.135.21 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.21->192.168.130.136.1050 TCP SYN ACK (seq: 8b892e46 ack: 12f799) ort 21 is seems to be OPEN. . 192.168.130.136.1050->192.168.130.135.21 TCP RST ACK (seq: 12f799 ack: 8b892e47)	192.168.128.15.1050->192.168.130.135.21 TCP SYN (seq: 16b89 ack: 0) 192.168.128.15.1051->192.168.130.135.21 TCP SYN (seq: 206fc ack: 0) 192.168.128.15.1052->192.168.130.135.21 TCP SYN (seq: 16fe4 ack: 0) 192.168.128.15.1053->192.168.130.135.21 TCP SYN (seq: 23ca8 ack: 0) 192.168.128.15.1053->192.168.130.135.21 TCP SYN (seq: 32ca8 ack: 0)
. 192.168.130.136.1050->192.168.130.135.79 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.79->192.168.130.136.1050 TCP RST ACK (seq: 0 ack: 12f799) ort 79 is seems to be CLOSED. . 192.168.130.136.1050->192.168.130.135.79 TCP RST ACK (seq: 12f799 ack: 1)	192.168.128.15.1055->192.168.130.135.21 TCP SYN (seq: 195e6 ack: 0) 192.168.128.15.1055->192.168.130.135.21 TCP SYN (seq: 26f2a ack: 0) 192.168.128.15.1057->192.168.130.135.21 TCP SYN (seq: 121dd ack: 0) 192.168.128.15.1059->192.168.130.135.21 TCP SYN (seq: 25d0 ack: 0) 192.168.128.15.1059->192.168.130.135.21 TCP SYN (seq: 25d0 ack: 0)
. 192.168.130.136.1050->192.168.130.135.80 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.80->192.168.130.136.1050 TCP SYN ACK (seq: 8b919779 ack: 12f799) ort 80 is seems to be OFEN. . 192.168.130.136.1050->192.168.130.135.80 TCP RST ACK (seq: 12f799 ack: 8b91977a)	192.168.128.15.1060->192.168.130.135.21 TCP SYN (seq: 94a7 ack: 0) 192.168.128.15.1061->192.168.130.135.21 TCP SYN (seq: 235af ack: 0) 192.168.128.15.1062->192.168.130.135.21 TCP SYN (seq: 17bh5 ack: 0) 192.168.128.15.1063->192.168.130.135.21 TCP SYN (seq: 17bh5 ack: 0) 192.168.128.15.1063->192.168.130.135.21 TCP SYN (seq: 20ef4 ack: 0) 192.168.128.15.1063->192.168.130.135.21 TCP SYN (seq: 20ef4 ack: 0)
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arting scanports v.1.0. TCP scanning by using SYN messages. tackID: SX	192.168.128.15.1070->192.168.130.135.21 TCP SYN (seq: 24e8c ack: 0) 192.168.128.15.1071->192.168.130.135.21 TCP SYN (seq: 21d63 ack: 0) 192.168.128.15.1072->192.168.130.135.21 TCP SYN (seq: 15fae ack: 0)
lected device: Realtek 8139-series PCI NIC	192.168.128.15.1073->192.168.130.135.21 TCP SYN (seq: 18088 ack: 0) 192.168.128.15.1074->192.168.130.135.21 TCP SYN (seq: 1ca25 ack: 0)
. 192.168.130.136.1050->192.168.130.135.21 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.21->192.168.130.136.1050 TCP SYN ACK (seq: 8ba2e74d ack: 12f799) ort 21 is seems to be OPEN. . 192.168.130.136.1050->192.168.130.135.21 TCP RST ACK (seq: 12f799 ack: 8ba2e74e>	192.168.128.15.1075->192.168.130.135.21 TCP SYN (seg: 1fe82 ack: 0) 192.168.128.15.1075->192.168.130.135.21 TCP SYN (seg: 2cbf ack: 0) 192.168.128.15.1077->192.168.130.135.21 TCP SYN (seg: 20332 ack: 0) 192.168.128.15.1077->192.168.130.135.21 TCP SYN (seg: 52c6 ack: 0) 192.168.128.15.1077->192.168.130.135.21 TCP SYN (seg: 52c6 ack: 0)
. 192.168.130.136.1050->192.168.130.135.79 TCP SYN (seq: 12f798 ack: 0) .192.168.130.135.79->192.168.130.136.1050 TCP RST ACK (seq: 0 ack: 12f799) .rt 79 is seems to be CLOSED. .192.168.130.136.1050->192.168.130.135.79 TCP RST ACK (seq: 12f799 ack: 1)	192.168.128.15.1080->192.168.130.135.21 TCP SYN (seq: 266d3 ack: 0) 192.168.128.15.1081->192.168.130.135.21 TCP SYN (seq: d165 ack: 0) 192.168.128.15.1082->192.168.130.135.21 TCP SYN (seq: 352a ack: 0) 192.168.128.15.1083->192.168.130.135.21 TCP SYN (seq: 1530b ack: 0) 192.168.128.15.1083->192.168.130.135.21 TCP SYN (seq: 162 ack: 0) 192.168.128.15.1083->192.168.130.135.21 TCP SYN (seq: 162 ack: 0) 192.168.128.15.1085->192.168.130.135.21 TCP SYN (seq: 162 ack: 0) 192.168.128.15.1085->192.168.130.135.21 TCP SYN (seq: 168 ack: 0)
. 192.168.130.136.1050->192.168.130.135.80 TCP SYN (seq: 12f798 ack: 0) . 192.168.130.135.80->192.168.130.136.1050 TCP SYN ACK (seq: 8bab55f7 ack: 12f799) ort 80 is seems to be OPEN. . 192.168.130.136.1050->192.168.130.135.80 TCP RST ACK (seq: 12f799 ack: 8bab55f8)	192.168.128.15.10085->192.168.130.135.21 TCP SVN (seq: 1a87e ack: 0) 192.168.128.15.10086->192.168.130.135.21 TCP SVN (seq: 1e50f ack: 0) 192.168.128.15.10087->192.168.130.135.21 TCP SVN (seq: 1e50f ack: 0) 192.168.128.15.10087->192.168.130.135.21 TCP SVN (seq: 1e50f ack: 0) 192.168.128.15.10087->192.168.130.135.21 TCP SVN (seq: 1e50f ack: 0) 192.168.128.15.10089->192.168.130.135.21 TCP SVN (seq: 1e50f ack: 0)



Shortcut to PORTAL.BAT		_0_	
G:\MAS4\Generic_agent\Server>cd	Specify the Attack		
G:\MAS4\Generic_agent>set PATH=	_ Intention		
NUTCROOT\bin;C:\PROGRA~1\RATION NUTCROOT\mksnt;C:\Program Files	N Name Description		
s\Microsoft Visual Studio\Commo on\Tools;C:\Program Files\Micros	1 IH Identification of Hosts		
.3.1\bin;C:\Program Files\Ratio	2 IS Identification of Services 3 IO Identification of Operating system		
am Files\Rational\Rational Test t Visual Studio\Common\Tools\Wiu	4 RE Shared Resource Enumeration		
\Bin;C:\Program Files\Microsoft al Studio\UC98\bin	5 UE Users and groups Enumeration 6 ABE Applications and Banners Enumeration		
	7 GAR Getting Access to Resources of the host		
G:\MAS4\Generic_agent>set PATH b;C:\PROGRA~1\RATIONAL\RATION~1	8 EP Escalating Privilege with regard t 9 CVR Confidentiality Violation Realization		
1;C:\PROGRA~1\RATIONAL\RATION~1` mmon\Tools\WinNT;C:\Program File	10 IVR Integrity Violation Realization		
les\Microsoft Visual Studio\Com	10 AVH Availability violation healization		
;C:\Program Files\JavaSoft\JRE\ Rational\Rose\TopLink\;C:\Progr			
SS\Bin;C:\Program Files\Microso oft Visual Studio\Common\MSDev9	Hacker Configuration		
;C:\Program Files\Microsoft Vis	Real IP-address	Save preceding attack realization	
G:\MAS4\Generic_agent>G:\MAS4\G		Generate attacks on net protocol level	
G:\MAS4\Generic_agent>javaDjav	Spoofed IP-address		rmation
G:\MAS4\Generic_agent)java -Djav b\jlex.jar;JavaLib\javacup jar	a Generat	ion Demons	
			rmation
	 Known Information about attacked Networks Networks 	- Hosts	irmation to new agent
	Net Name Net IP	Host Name Host IP	ainHak to MainHak
G:\MAS4\Generic_agent\Server>cd	AIL 192.168.130.0	192.168.130.138	
G:\MAS4\Generic_agent>set PATH=G \RATION~1\NUTCROOT\bin;C:\PROGRA	· · · · · · · · · · · · · · · · · · ·	192.168.130.139 192.168.130.140	
1\RATIONAL\RATION~1\NUTCROOT\mks		192.168.130.141	🖓 🤃 Portal Component. 💦 🗔 🔯
on\Tools\WinNT;C:\Program Files\ ogram Files\Microsoft Visual Stu		192.168.130.135	
al Studio/UC98/hin:C:/Program Fi			
onal\common;C:\Program Files\Rat tional Test;C:\j2sdk1.4.1_02\bin tudio\Common\Tools\WinNT;C:\Prog			
\Bin;C:\Program Files\Microsoft			nHak _ 💷 🗡
osoft Visual Studio\VC98\bin	Define Known Information	Show All Hosts	
G:∖MAS4\Generic_agent>java −cla AILLogger	Object of Attack		es
Thogger	Object of Attack		
Advanced			
Intention: IH IP-address: 192.168.130.138 QK Cancel			
		68.130.138 OK Cancel	
	THE REPORT OF THE OWNER		

Parameters of attack realization outcome

- NS (Number of attack Steps) – number of terminal level attack actions;

- **PIR (Percentage of Intention Realization)** – percentage of the hacker's intentions realized successfully (for "Reconnaissance" it is a percentage of objects about which the information has been received; for "Implantation and threat realization" it is a percentage of successful realizations of the common attack goal on all runs);

- PAR Percentage of Attack actions Realization – percentage of "positive" messages (responses) of the Network Agent on attack actions (the "positive" messages are designated in attack visualization window by green lines);

- **PFB (Percentage of Firewall Blockage)** – percentage of attack actions blockage by firewall (red lines in attack visualization window);

- PRA (Percentage of Reply Absence) - percentage of "negative" messages (responses) of the Network Agent on attack actions (gray lines in attack visualization window).

Example of experiment results for intention "Gaining Access to host Resources"



Configurations of firewalls: 1 - Both Net & Personal firewalls are active; 2 - Only Net firewall is active; 3 - Only Personal firewall is active; 4 - None of firewalls is active



Modeling and simulation of cyberwarfare (between malefactors' teams and securty teams)

Research objectives

Development of the formal framework, models, architecture, and software for agent-based modeling and simulation of adversarial interaction of teams of malefactors and security teams aimed to create theoretical bases for construction of *integrated* intrusion-aware trusted security systems operating in adversarial environments.

Interaction of team of malefactors and computer network assurance system components





Abstract model of team interaction





Main Classes of Attack and Defense Parameters. Parameters of Defense Efficiency

Attack module

- Victim type
- Attack type
- Impact on the victim
- Attack rate dynamics
- Persistent of agent set
- Possibility of exposure
- Source address validity
- Degree of automation
- Deployment location
- Mechanism of cooperation
- Covered defense stages
- Attack detection technique
- Attack source detection technique
- Attack prevention/counteraction technique
- Model data gathering technique
- Determination of deviation from model data

Efficiency Parameters:

- List of detectable attacks
- Volume of the input traffic before and after filters
- Percent of the normal traffic and the attack traffic on entrance to attacked network
- Rate of dropped legitimate traffic (false positive rate)
- Rate of admitted attack traffic (false positive rate)
- Attack detection and attack reaction times
- Computational complexity
- etc.

RE-TRUST Kick-off Workshop, September 18-19, 2006

Defense module

Architecture of Simulation Environment



User Interface of Simulation Environment





Learning Mode (1)

- The main task of learning mode is to create the model of generic traffic for the given network.
- *The clients* send the requests to the server and it replies.
- At this time *sampler* analyses requests and uses them to form the models and parameters for defense different methods.
- During the learning it is possible to watch the *change of traffic models*.

Learning Mode (2)



List of hosts that sent requests to server and <u>hops</u> to them after 300 sec of learning

Learning Mode (3)

General

class std::vector<AR NormIP *> {

many new addresses in the beginning

many new addresses in the interval between 0 and 50 seconds

AR NormIP * *(nipv.getVectorPtr())[0] = IP=10.0.0.33

AR_NormIP * *(nipv.getVectorPtr())[1] = IP=10.0.0.36

AR NormIP * *(nipv.getVectorPtr())[2] = IP=10.0.0.37

AR NormIP * *(nipv.getVectorPtr())[3] = IP=10.0.0.25

AR_NormIP * *(nipv.getVectorPtr())[4] = IP=10.0.0.27

AR_NormIP * *(nipv.getVectorPtr())[5] = IP=10.0.0.35 AR_NormIP * *(nipv.getVectorPtr())[6] = IP=10.0.0.24

AR_NormIP * *(nipv.getVectorPtr())[7] = IP=10.0.0.23

AR NormIP * *(nipv.getVectorPtr())[8] = IP=10.0.0.31

AR NormIP * *(nipv.getVectorPtr())[\$] = IP=10.0.0.28

AR NormIP * *(nipv.getVectorPtr())[1] = IP=10.0.0.29

AR_NormIP * *(nipv.getVectorPtr())[12] = IP=10.0.0.22

AR_NormIP * *(nipv.getVectorPtr())[13] = IP=10.0.0.30

AR NormIP * *(nipv.getVectorPtr())[14] = IP=10.0.0.32

(std::vector <AR_NormIP *>) ...p[0].ad_stats app.*(nipv....

Time=12.0008

Time=14.0003

Time=14.0001

Time=23.1377

Time=21.9945

Time=21.9947

Time=32.034

Time=35,1563

Time=33.8222

Time=37.3439

Time=39.925

Time=42.0896

Time=45.5916

Time=100.002



Change of <u>new IP</u> addresses amount

List of <u>clients</u> requested server and considered as <u>legitimate</u> after 300 sec of learning

AR_NormIP * *(nipv.getVectorPtr())[10] = IP=10.0.0.26 Time=39.8259

Learning Mode (4)

The maximum value was 1742.4 bit/s

Values of bits in interval 10 seconds



Change of BPS (bit per
second) parameterValues of transmitted bits for different
hosts

Decision Making and Acting (1)

- Normal work (interval 0 300 seconds)
- Defense team: Formation, start using BPS method
- Attack team: Formation

-<u>Attack team</u>: After 300 seconds - begins the attack actions (intensity of attack for every daemon - 0.5, **no IP spoofing**)

- <u>Defense team</u>: data processing, attack detecting (**using BPS**) and reacting (interval 300 – 350 seconds)

- <u>Defense team</u>: blocking the attack, destroying some attack agents (interval 300 – 600 seconds)

Decision Making and Acting (2)

- <u>Attack team</u>: After 600 seconds - **automatic adaptation** (redistributing the intensity of attack (0.83), changing the method of **IP spoofing (Random)**)

- <u>Defense team</u>: data processing, failing to detect the attack (using BPS method) – Detector sees that the input channel throughput has noticeably lowered, but does does not receive any anomaly report from sampler because BPS does not work.

- <u>Defense team</u>: Changing defense method on **SIPM** (automatic adaptation).

<u>Defense team</u>: data processing, attack detecting (using
 SIPM method) and reacting – (interval 600 – 700 seconds)



Cooperation between defense teams

Models of cooperation between distributed defense teams: (1) filter-level cooperation (2) sampler-level cooperation (3) "poor" cooperation: (4) "full" cooperation Such cooperation schemas are used in the cooperative DDoS defense methods: COSSACK, Perimeter-based DDoS defense, DefCOM, Gateway-based, ACC pushback, MbSQD, SOS, tIP router architecture, etc.)







• Security analysis of computer networks



Main features of the Approach

- Based on malefactor's action simulation and integrated family of various expert knowledge models
- Two main phases:
 - (1) construction of attack graph and
 - (2) computation of different security metrics using combination of qualitative techniques of risk analysis
- Taking into account diversity of malefactor's positions, intentions and experience;
- Estimating the influence of different configuration and policy data;
- Taking into account not only attack actions (which use vulnerabilities), but the common actions of legitimate users and reconnaissance actions;

Main features of the Approach (cont'd)

- Investigation of various threats for different resources;
- Detection of "weak" places;
- Usage of up to date databases of vulnerabilities (NVD, OSVDB);
- The "CVSS. Common Vulnerability Scoring System" approach is used for computation of a part of primary security metrics;
- Comparing calculated metrics and user requirements
- The qualitative techniques of risk analysis are used for computation of security metrics (in particular SANS/GIAC and FRAP technique).

Generalized Architecture



Generalized Attack Scenarios



Generating Common Attack Graph

- Realization of actions which are intended for malefactor's movement from one host onto another in the following cases:
 - There is a possibility to realize the attack actions which use vulnerabilities of software and hardware and which require that the malefactor already have the privileges of local user
 - Movement of the malefactor into the attacked host allows him to penetrate the another segment of network
 - Movement of the malefactor into attacked host allows him to use "trust relationships"
- Realization of reconnaissance actions for detection of "live" hosts (ex., "ping" utility)
- Realization of reconnaissance scenarios (the sets of actions) for each of the detected hosts (ex, "nmap OS", "nmap services", "banners")
- Realization of attack actions which use vulnerabilities of software and hardware, and common actions of ordinary users



Model of Security Level Evaluation

Model of Security Level Evaluation consists of:

- Security metrics
- Rules (formulas) for their calculation

Two approaches for security level evaluation:

- Qualitative express assessment (!)
- Quantitative computation

Taxonomy of Security Metrics (SM)

- According to division of objects of attack graph
 - SM of base objects (hosts, attack actions)
 - SM of complex objects (routes, threats, graph)
- According to the order of calculation
 - Primary (received directly from attack graph)
 - Secondary (calculated on the basis of primary)
- Whether metrics are used for evaluation of general security level
 - Basic (are used for evaluation of general security level)
 - Auxiliary (are not used for evaluation of general security level)

Basic Security Metrics

- Criticality Level of the host Criticality(h)
- Criticality Level of attack action Severity(a)
- Damage Level of attack action Mortality(a,h)
- Damage Level of route Mortality(S) or threat Mortality(T)
- Access Complexity Level AccessComplexity(a), AccessComplexity(S), AccessComplexity(T)
- Admissibility of threat realization Realization(T)
- Risk Level of threat RiskLevel(T)
- General Security Level of network SecurityLevel

Technique of General Security Level Evaluation

- Calculation of security metrics of basic and complex objects
- Estimation of qualitative risk level for all threats
- Evaluation of security level of analyzed computer network on basis of received values of risk levels for threats
Implementation: User Interface (1)



Implementation: User Interface (2)





Intrusion detection

Different Sources of information



Multi-agent Architecture of Raw Data Preprocessing and Intrusion Detection



Case Study: Anomaly Detection in Computer Network

Computer security status: {Normal, Abnormal}.

- Types of attacks constituting class "*Abnormal*": {Probing, Remote to local (R2L); Denial of service (DOS) and User to root (U2R)".
- Instances of attacks of respective classes: {SYN-scan, FTP-crack attack, SYN flood, and PipeUpAdmin}.
- Information Source: Network traffic raw data.
- Data source 1: Stream of binary vectors specifying stream of headers of *IP* packets within a connection (sequence of binary vectors).
- Data source 2: Statistical attributes of particular connections manifesting in input traffic. (duration, status, total number of connection packets and also other attributes specifying statistics of connections).
- Data source 4: Statistical attributes of traffic for long time intervals (composed of the same statistics as previous ones averaged over 100 connections).



Architecture of host-based MIDS components



1,2,3 – levels of processing; 1,2,...,7 – types of agents.











• Deception systems, honeynets

Deception System (DS) network architecture







Security policy specification and checking





Examples of Verification Modules

Event Calculus verification module

- security policies and system description are translated into domaindependent Event Calculus axiomatic
- conflict predicates are introduced
- abductive inference is used for conflict detection
- Implemented in Jess

SPIN verification module

- security policies and system description are translated into Promela data structures, processes, and assertions
- policy conflicts introduced as additional assertions
- impemented in SPIN-Promela

Implementation of security policy verification system

- VerificationManager
 - registers, loads and authenticates verification modules
 - invokes verify() method of verification modules for policy consistency and enforceability checks
 - debugs policy if conflicts have been found
 - open for new verification modules
 - semi-lattice-based, semantic approach
- VerificationModule
 - checks for policy consistency and predicts policy violations
 - resolves conflicts by means of a resolution strategy
- permits independent development of module
 - speed up implementation
- needs administrator to check or change results
 - choosing between proposed alternatives

Model Checking VM: functional model (IDEF)



Model Checking VM: common architecture



Event Calculus VM: functional model (IDEF)



Event Calculus VM: technique



Security system and corresponding events representation in EC axiomatics

Policy conflict formalization as abductive queries

Interpreting assertion violation as conflicts

Security policy verification system GUI

SEC: security policy checker		
le Modules Help		**************************************
	Select verification module(s) for Registered modules	++++++++++++++++++++++++++++++++++++++
Event Calculus verification module		View In initially_false(F) iff [[F=authorization_allowed(User.
SPIN verification module		<pre>write, server)], [F=authorization_denied(User, write, server)]]. View linitiates(E,F,T) iff [[subject_role</pre>
Test verification module		(User, fTPWriteAllowedRole), E=request_authorization View k (User, write, server), F=authorization_allowed(User,
-		write, server)], [subject role
Test verification module 2		View I (User, fTPWriteDeniedRole), E=request_authorization (User, write, server), F=authorization_denied(User, write, server)]].
elect policy and system description	<pre>/2_demo_authorization_conflict</pre>	
Verification Results		ciff('res/sec/ec/policies.alp', [holds_at (authorization allowed(User, Action, Target), T).
VERIFICATION RESULT		<pre>(authorization_allowed(User, Action, Target), T), holds_at(authorization_denied(User, Action, Target), T)], A).</pre>
System Y2_dem	no_authorization_conflict.xml	Sector sector en esta esta esta esta esta esta Sector esta esta esta esta esta esta esta esta
Module name	SPIN verification module	CONFLICT(S) FOUND [subject_role(_107066_fTPWciteAllowedRole)
Result		[subject_role(_107066,fTPWriteAllowedRole), subject_role(_107066,fTPWriteDeniedRole), happens (request_authorization(_107066,write,server),_17027), happens(request_authorization
Conflict description		FTPWriteAllowedRole
A	B	conflicts.roles[1] =
	IN verification module Intradictory rules:FTPWriteAllo	Strategy FTPWriteDeniedRole userRoles [0] = 1
	solvable	Resolve userRoles[1] = 1 31: proc 1 (RoleAssigner) line 31 "model"
Strategies DTF	P, Manual modificatio	Next conflict (state 5) <valid end="" state=""></valid>
Deactivate or edit rules	× Select resolution strat	31: proc 0 (:init:) line 51 "model" (state 7) <valid end="" state=""></valid>
Rules	Available strategies	('never') line 68 "model"
Deactivate FTPWriteAllowedRule	DTP	orm 3 processes created
Deactivate FTPWriteDeniedRule	Manual modification	******
		CONFLICT(S) FOUND Rule-1: FTPWriteAllowedRule
	Select	Rule-2: FTPWriteDeniedRule
Apply Close	301000 0	incel i i i i i i i i i i i i i i i i i i i

Thanks!



For more information and related publications please contact

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